

1) Trovare il codice assembly MIPS corrispondente al seguente programma (utilizzando solo e unicamente istruzioni dalla tabella sottostante), rispettando le convenzioni di utilizzazione dei registri dell'assembly MIPS (riportate in calce, per riferimento).

```
char buff[80] = "Britney Spears\n";

char to_upper(char c)
{
    if (c >= 'a' && c <= 'z') c -= 0x20;
    return (c);
}

char *myfun(int n, char *p, char c, float f, double d)
{
    char *r, *s = p;

    while (*s++ != '\0' && n-- > 0) {
        *s = to_upper(*s);
        if (*s == c) r = s;
    }

    if (f * d < 0) {
        f = -f;
        r = myfun(7, r, c, f, d);
    }
    return (r);
}

main()
{
    char *p;

    printf(buff);
    p = myfun(7, buff, 'E', 1, -1);
    printf(p);
}
```

MIPS instructions

Instruction	Example	Meaning	Comments
add	add \$1,\$2,\$3	\$1 = \$2 + \$3	3 operands; exception possible
subtract	sub \$1,\$2,\$3	\$1 = \$2 - \$3	3 operands; exception possible
add immediate	addi \$1,\$2,100	\$1 = \$2 + 100	+ constant; exception possible
subtract immediate	subi \$1,\$2,100	\$1 = \$2 - 100	- constant; exception possible
multiplication	mult \$1, \$2	Hi,Lo = \$1 x \$2	64-bit Signed Product ; result in Hi,Lo
division	div \$1, \$2	Hi = \$1 % \$2, Lo = \$1 / \$2	Signed division
move from Hi	mhi \$1	\$1 = Hi	Create copy of Hi
move from Lo	mlo \$1	\$1 = Lo	Create copy of Lo
and	and \$1,\$2,\$3	\$1 = \$2 & \$3	3 register operands; Logical AND
or	or \$1,\$2,\$3	\$1 = \$2 \$3	3 register operands; Logical OR
nor	nor \$1,\$2,\$3	\$1 = !((\$2 \$3))	3 register operands; Logical NOR
xor	xor \$1,\$2,\$3	\$1 = \$2 ^ \$3	3 register operands; Logical XOR
and immediate	andi \$1,\$2,100	\$1 = \$2 & 100	Logical AND register, constant
or immediate	ori \$1,\$2,100	\$1 = \$2 100	Logical OR register, constant
xor immediate	xori \$1,\$2,100	\$1 = \$2 ^ 100	Logical XOR register, constant
shift left logical	sll \$1,\$2,10	\$1 = \$2 << 10	Shift left by constant
shift right logical	srl \$1,\$2,10	\$1 = \$2 >> 10	Shift right by constant
load word	lw \$1,100(\$2)	\$1 = Memory[\$2+100]	Data from memory to register
load byte	lb \$1,100(\$2)	\$1 = Memory[\$2+100]	Data from memory to register
load byte unsigned	lbu \$1,100(\$2)	\$1 = Memory[\$2+100]	Data from mem. to reg.: no sign extension
store word	sw \$1,100(\$2)	Memory[\$2+100] = \$1	Data from register to memory
store byte	sb \$1,100(\$2)	Memory[\$2+100] = \$1	Data from register to memory
load address	la \$1,var	\$1 = &var	Load variable address
branch on equal	beq \$1,\$2,100	if (\$1 == \$2) go to PC+4+100	Equal test; PC relative branch
branch on not equal	bne \$1,\$2,100	if (\$1 != \$2) go to PC+4+100	Not equal test; PC relative
set on less than	slt \$1,\$2,\$3	if (\$2 < \$3) \$1 = 1; else \$1 = 0	Compare less than; 2's complement
set on less than immediate	slti \$1,\$2,100	if (\$2 < 100) \$1 = 1; else \$1 = 0	Compare < constant; 2's complement
set on less than unsigned	sltu \$1,\$2,\$3	if (\$2 < \$3) \$1 = 1; else \$1 = 0	Compare less than; natural number
set on less than imm. unsigned	sltiu \$1,\$2,100	if (\$2 < 100) \$1 = 1; else \$1 = 0	Compare constant; natural number
jump	j 10000	go to 10000	Jump to target address
jump register	jr \$31	go to \$31	For switch, procedure return
jump and link	jal 10000	\$31 = PC + 4; go to 10000	For procedure call
add.s add.d	add.s \$f0,\$f2,\$f4	\$f0 = \$f2 + \$f4	Single and double precision add
sub.s sub.d	add.s \$f0,\$f2,\$f4	\$f0 = \$f2 - \$f4	Single and double precision subtraction
mul.s mul.d	mul.s \$f0,\$f2,\$f4	\$f0 = \$f2 * \$f4	Single and double precision multiplication
div.s div.d	div.s \$f0,\$f2,\$f4	\$f0 = \$f2 / \$f4	Single and double precision division
mov.s mov.d	mov.s \$f0,\$f2	\$f0 ← \$f2	Single and double precision move
abs.s abs.d	abs.s \$f0,\$f2	\$f0 = ABS(\$f2)	Single and double precision absolute value
c.lt.s c.lt.d (eq.ne.lt.gt.ge)	c.lt.s \$f0,\$f2	Temp = (\$f0 < \$f2)	Single and double: compare \$f0 and \$f2 <= <=> >=>
branch on false	bnef label	If (Temp == false) go to label	Temp is 'Condition-Code'
branch on true	bnef label	If (Temp == true) go to label	Temp is 'Condition-Code'
load floating point (32bit)	lwc1 \$f0,0(\$1)	\$f0 ← Memory[\$1]	
store floating point (32bit)	swc1 \$f0,0(\$1)	Memory[\$1] ← \$f0	
convert single into double	cvt.d.s \$f0,\$f2	\$f0 = (double)\$f2	Also cvt.s.d (viceversa)
convert single into integer	cvt.w.s \$t0,\$f0	\$t0 = (int)\$f0	Also cvt.s.w (viceversa)

Register Usage

Name	Register Num.	Usage
\$zero	0	The constant value 0
\$s0-\$s7	16-23	Saved
\$t0-\$t9	8-15,24-25	Temporaries
\$a0-\$a3	4-7	Arguments

Name	Register Num.	Usage
\$v0-\$v1	2-3	Results
\$fp, \$sp	30,29	Frame pointer, stack pointer
\$ra, \$gp	31,28	return address, global pointer
\$k0-\$k1	26,27	Kernel usage

Name	Usage
\$f0, \$f1, ..., \$f31	Single precision floating point registers
\$f0, \$f2, ..., \$f30	Double precision floating point registers

Si presenta una possibile soluzione:

```
.data
buff: .asciiz "Britney Spears\n"
      .space 64
zerof: .float 0.0
punof: .float 1.0
munod: .double -1.0
tempf: .float 0.0

.text
.globl main

to_upper:
    slti    $t0, $a0, 'a'
    bne    $t0, $0, ret_to_upper
    addi   $t1, $0, 'z'
    slt    $t2, $t1, $a0
    bne    $t0, $0, ret_to_upper
    addi   $a0, $a0, -0x20
ret_to_upper:
    add    $v0, $a0, $0
    jr     $ra

myfun:
    addi   $sp, $sp, -32 # 8 per var.locali, 4 per $ra, 4 $fp, 16 arg.
    sw     $fp, 12($sp)
    add    $fp, $sp, $0 # inizializza il nuovo frame pointer
                    # d -->inserito dal chiamante
    sw     $a3, 28($fp) # f
    sw     $a2, 24($fp) # c
    sw     $a1, 20($fp) # p
    sw     $a0, 16($fp) # n
    sw     $ra, 8($fp) # ind. ritorno
    sw     $s1, 4($fp) # $s1 = s
    sw     $s0, 0($fp) # $s0 = r

    add    $s1, $a1, $0 # s = p
while_start:
    lb     $t0, 0($s1) # *s
    addi   $s1, $s1, 1 # s++
    beq    $t0, $0, while_end # salta se la prima cond. e' falsa
    add    $t1, $a0, $0 # n
    addi   $a0, $a0, -1 # n--
    slt    $t2, $0, $t1 # 0 <? n
    beq    $t2, $0, while_end # salta se la seconda cond. e' falsa

    sw     $a0, 16($fp) # save $a0, richiesto da to_upper
    lb     $t0, 0($s1) # nuovo *s
    add    $a0, $t0, $0 # setup del parametro di input
    jal    to_upper # chiama to_upper
    lw     $a0, 16($fp) # ripristina $a0
    sb     $v0, 0($s1) # memorizza risultato, *s = ...
    lb     $t0, 0($s1) # nuovo *s
    bne    $t0, $a2, while_start

    add    $s0, $s1, $0 # r = s
    j     while_start

while_end:
    lwcl   $f0, 28($fp) # carica f in $f0
    lwcl   $f2, 32($fp) # carica d in ($f2,$f3)
    lwcl   $f3, 36($fp) # 64 bit...
    cvt.d.s $f4, $f0 # converti f in double
    la     $t0, zerof
```

```

    lwc1 $f30, 0($t0)
    cvt.d.s $f8, $f30 # ($f8,$f9) = 0.0 (double)

    mul.d $f6, $f2, $f4 # (double)f * d
    c.lt.d $f6, $f8 # f*d <? 0
    bclf if_end

    sub.s $f12, $f30, $f12 # f = -f
    addi $a0, $0, 7 # $a0 = 7
    add $a1, $s0, $0 # $a1 = r
    # $a2 e' a posto = c
    la $t0, tempf # espediente per non usare
    swc1 $f12, 0($t0) # la mfc1 $a3, $f12
    lw $a3, 0($t0) # cosi' sistemmo $a3

    addi $sp, $sp, -8
    swc1 $f2, 0($sp)
    swc1 $f3, 4($sp) # quinto parametro = d

    jal myfun

    add $s0, $v0, $0 # r = risultato di myfun

if_end:
    add $v0, $s0, $0 # parametro restituito

    lw $a0, 16($fp) # ripristina a0, a1, a2, a3
    lw $a1, 20($fp) # (necessario a causa della chiamata ricorsiva)
    lw $a2, 24($fp) #
    lw $a3, 28($fp) #
    lw $ra, 8($fp) # ripristina $ra, $fp
    lw $fp, 12($fp) # (necessario a causa della chiamata ricorsiva)
    addi $sp, $sp, 40 # ripristina lo stack e butta i parametri locali r e s
    jr $ra

main:
    la $a0, buff
    addi $v0, $0, 4
    syscall

    addi $a0, $0, 7
    la $a1, buff
    addi $a2, $0, 'E'
    la $t0, punof
    lw $a3, 0($t0)
    addi $sp, $sp, -8
    la $t0, munod
    lw $t1, 0($t0)
    sw $t1, 0($sp)
    lw $t2, 4($t0)
    sw $t2, 4($sp)

    jal myfun

    add $a0, $v0, $0
    addi $v0, $0, 4
    syscall

    addi $v0, $0, 10
    syscall

```