

1) Trovare il codice assembly MIPS corrispondente al seguente programma (utilizzando solo e unicamente istruzioni dalla tabella sottostante), rispettando le convenzioni di utilizzazione dei registri dell'assembly MIPS (riportate in calce, per riferimento).

```

char buff[80] = "Britney Spears\n";

char to_upper(char c)
{
    if (c >= 'a' && c <= 'z') c -= 0x20;
    return (c);
}

char *myfun(int n, char *p, char c, float f, double d)
{
    char *r, *s = p;

    while (*s++ != '\0' && n-- > 0) {
        *s = to_upper(*s);
        if (*s == c) r = s;
    }

    if (f * d < 0) {
        f = -f;
        r = myfun(7, r, c, f, d);
    }
    return (r);
}

main()
{
    char *p;

    printf(buff);
    p = myfun(7, buff, 'E', 1, -1);
    printf(p);
}
    
```

**MIPS instructions**

Instruction	Example	Meaning	Comments
add	add \$1,\$2,\$3	\$1 = \$2 + \$3	3 operands; exception possible
subtract	sub \$1,\$2,\$3	\$1 = \$2 - \$3	3 operands; exception possible
add immediate	addi \$1,\$2,100	\$1 = \$2 + 100	+ constant; exception possible
subtract immediate	subi \$1,\$2,100	\$1 = \$2 - 100	- constant; exception possible
multiplication	mult \$1, \$2	Hi,Lo = \$1 x \$2	64-bit Signed Product ; result in Hi,Lo
division	div \$1, \$2	Hi = \$1 % \$2, Lo = \$1 / \$2	Signed division
move from Hi	mhi \$1	\$1 = Hi	Create copy of Hi
move from Lo	mlo \$1	\$1 = Lo	Create copy of Lo
and	and \$1,\$2,\$3	\$1 = \$2 & \$3	3 register operands; Logical AND
or	or \$1,\$2,\$3	\$1 = \$2   \$3	3 register operands; Logical OR
nor	nor \$1,\$2,\$3	\$1 = !((\$2   \$3))	3 register operands; Logical NOR
xor	xor \$1,\$2,\$3	\$1 = \$2 ^ \$3	3 register operands; Logical XOR
and immediate	andi \$1,\$2,100	\$1 = \$2 & 100	Logical AND register, constant
or immediate	ori \$1,\$2,100	\$1 = \$2   100	Logical OR register, constant
xor immediate	xori \$1,\$2,100	\$1 = \$2 ^ 100	Logical XOR register, constant
shift left logical	sll \$1,\$2,10	\$1 = \$2 << 10	Shift left by constant
shift right logical	srl \$1,\$2,10	\$1 = \$2 >> 10	Shift right by constant
load word	lw \$1,100(\$2)	\$1 = Memory[\$2+100]	Data from memory to register
load byte	lb \$1,100(\$2)	\$1 = Memory[\$2+100]	Data from memory to register
load byte unsigned	lbu \$1,100(\$2)	\$1 = Memory[\$2+100]	Data from mem. to reg.: no sign extension
store word	sw \$1,100(\$2)	Memory[\$2+100] = \$1	Data from register to memory
store byte	sb \$1,100(\$2)	Memory[\$2+100] = \$1	Data from register to memory
load address	la \$1,var	\$1 = &var	Load variable address
branch on equal	beq \$1,\$2,100	if (\$1 == \$2) go to PC+4+100	Equal test; PC relative branch
branch on not equal	bne \$1,\$2,100	if (\$1 != \$2) go to PC+4+100	Not equal test; PC relative
set on less than	slt \$1,\$2,\$3	if (\$2 < \$3) \$1 = 1; else \$1 = 0	Compare less than; 2's complement
set on less than immediate	slti \$1,\$2,100	if (\$2 < 100) \$1 = 1; else \$1 = 0	Compare < constant; 2's complement
set on less than unsigned	sltu \$1,\$2,\$3	if (\$2 < \$3) \$1 = 1; else \$1 = 0	Compare less than; natural number
set on less than imm. unsigned	sltiu \$1,\$2,100	if (\$2 < 100) \$1 = 1; else \$1 = 0	Compare constant; natural number
jump	j 10000	go to 10000	Jump to target address
jump register	jr \$31	go to \$31	For switch, procedure return
jump and link	jal 10000	\$31 = PC + 4; go to 10000	For procedure call
add.s add.d	add.s \$f0,\$f2,\$f4	\$f0 = \$f2 + \$f4	Single and double precision add
sub.s sub.d	add.s \$f0,\$f2,\$f4	\$f0 = \$f2 - \$f4	Single and double precision subtraction
mul.s mul.d	mul.s \$f0,\$f2,\$f4	\$f0 = \$f2 * \$f4	Single and double precision multiplication
div.s div.d	div.s \$f0,\$f2,\$f4	\$f0 = \$f2 / \$f4	Single and double precision division
mov.s mov.d	mov.s \$f0,\$f2	\$f0 ← \$f2	Single and double precision move
abs.s abs.d	abs.s \$f0,\$f2	\$f0 = ABS(\$f2)	Single and double precision absolute value
c.lt.s c.lt.d (eq.ne.lt.gt.ge)	c.lt.s \$f0,\$f2	Temp = (\$f0 < \$f2)	Single and double: compare \$f0 and \$f2 <= <=> >=>
branch on false	bnef label	If (Temp == false) go to label	Temp is 'Condition-Code'
branch on true	bnef label	If (Temp == true) go to label	Temp is 'Condition-Code'
load floating point (32bit)	lwc1 \$f0,0(\$1)	\$f0 ← Memory[\$1]	
store floating point (32bit)	swc1 \$f0,0(\$1)	Memory[\$1] ← \$f0	
convert single into double	cvt.d.s \$f0,\$f2	\$f0 = (double)\$f2	Also cvt.s.d (viceversa)
convert single into integer	cvt.w.s \$t0,\$f0	\$t0 = (int)\$f0	Also cvt.s.w (viceversa)

**Register Usage**

Name	Register Num.	Usage
\$zero	0	The constant value 0
\$s0-\$s7	16-23	Saved
\$t0-\$t9	8-15,24-25	Temporaries
\$a0-\$a3	4-7	Arguments

Name	Register Num.	Usage
\$v0-\$v1	2-3	Results
\$fp, \$sp	30,29	Frame pointer, stack pointer
\$ra, \$gp	31,28	return address, global pointer
\$k0-\$k1	26,27	Kernel usage

Name	Usage
\$f0, \$f1, ..., \$f31	Single precision floating point registers
\$f0, \$f2, ..., \$f30	Double precision floating point registers