

1) [5/40] Ricavare la rappresentazione binaria in formato IEEE-754 singola precisione del numero 0.045 supponendo di effettuare un arrotondamento al piu' vicino valore in virgola mobile rappresentabile nel suddetto formato.

2) [35/40] Trovare il codice assembly MIPS corrispondente del seguente programma (**utilizzando solo e unicamente istruzioni dalla tabella sottostante**), **rispettando le convenzioni di utilizzazione dei registri dell'assembly MIPS** (riportate in calce, per riferimento). La funzione exp() riceve un numero di tipo "float" e ne restituisce l'esponenziale (e' una funzione esterna da dichiarare opportunamente, in particolare \$v0=exp(\$a0)...).

```
int num_cond(float T[3][], int n, float *nc)
{
    int r = 0, j=1;
    int k;

    while (j <= n) {
        for (k = 0; k < n; ++k) {
            if (T[j-1][k] != 0) {
                *nc += 1 / exp(-T[j-1][k]);
            } else {
                r = 1;
            }
        }
        ++j;
    }
    return (r);
}

float f = 0.0;
float A[][] =
{{1.0,2.0,3.0},{4.0,5.0,6.0},{7.0,8.0,9.0}};

main()
{
    int ec;

    ec = num_cond(A, 3, &f);

    printf("esito=");
    printf(ec);
    printf(" n.cond=");
    printf(f);
    printf("\n");
}
```

**MIPS instructions**

Instruction	Example	Meaning	Comments
add	add \$1,\$2,\$3	\$1 = \$2 + \$3	3 operands; exception possible
subtract	sub \$1,\$2,\$3	\$1 = \$2 - \$3	3 operands; exception possible
add immediate	addi \$1,\$2,100	\$1 = \$2 + 100	+ constant; exception possible
subtract immediate	subi \$1,\$2,100	\$1 = \$2 - 100	- constant; exception possible
multiplication	mult \$1, \$2	Hi,Lo= \$1 x \$2	64-bit Signed Product ; result in Hi,Lo
division	div \$1, \$2	Hi= \$1 % \$2, Lo = \$1 / \$2	Signed division
move from Hi	mfhi \$1	\$1 = Hi	Create copy of Hi
move from Lo	mflo \$1	\$1 = Lo	Create copy of Lo
and	and \$1,\$2,\$3	\$1 = \$2 & \$3	3 register operands; Logical AND
or	or \$1,\$2,\$3	\$1 = \$2   \$3	3 register operands; Logical OR
nor	nor \$1,\$2,\$3	\$1 = ~( \$2   \$3)	3 register operands; Logical NOR
xor	xor \$1,\$2,\$3	\$1 = \$2 ^ \$3	3 register operands; Logical XOR
and immediate	andi \$1,\$2,100	\$1 = \$2 & 100	Logical AND register, constant
or immediate	ori \$1,\$2,100	\$1 = \$2   100	Logical OR register, constant
xor immediate	xori \$1,\$2,100	\$1 = \$2 ^ 100	Logical XOR register, constant
shift left logical	sll \$1,\$2,10	\$1 = \$2 << 10	Shift left by constant
shift right logical	srl \$1,\$2,10	\$1 = \$2 >> 10	Shift right by constant
load word	lw \$1,100(\$2)	\$1 = Memory[\$2+100]	Data from memory to register
load byte	lb \$1,100(\$2)	\$1 = Memory[\$2+100]	Data from memory to register
load byte unsigned	lbu \$1,100(\$2)	\$1 = Memory[\$2+100]	Data from mem. to reg.; no sign extension
store word	sw \$1,100(\$2)	Memory[\$2+100] = \$1	Data from register to memory
store byte	sb \$1,100(\$2)	Memory[\$2+100] = \$1	Data from register to memory
load address	la \$1,var	\$1 = &var	Load variable address
branch on equal	beq \$1,\$2,100	if (\$1 == \$2) go to PC+4+100	Equal test; PC relative branch
branch on not equal	bne \$1,\$2,100	if (\$1 != \$2) go to PC+4+100	Not equal test; PC relative
set on less than	slt \$1,\$2,\$3	if (\$2 < \$3) \$1 = 1; else \$1 = 0	Compare less than; 2's complement
set on less than immediate	slti \$1,\$2,100	if (\$2 < 100) \$1 = 1; else \$1 = 0	Compare < constant; 2's complement
set on less than unsigned	sltu \$1,\$2,\$3	if (\$2 < \$3) \$1 = 1; else \$1 = 0	Compare less than; natural number
set on less than imm. unsigned	sltiu \$1,\$2,100	if (\$2 < 100) \$1 = 1; else \$1 = 0	Compare constant; natural number
jump	j 10000	go to 10000	Jump to target address
jump register	jr \$31	go to \$31	For switch, procedure return
jump and link	jal 10000	\$31 = PC + 4; go to 10000	For procedure call
add.s add.d	add.x \$F0,\$F2,\$F4	\$F0=\$F2+\$F4	Single and double precision add
sub.s sub.d	add.x \$F0,\$F2,\$F4	\$F0=\$F2-\$F4	Single and double precision subtraction
mul.s mul.d	mul.x \$F0,\$F2,\$F4	\$F0=\$F2*\$F4	Single and double precision multiplication
div.s div.d	div.x \$F0,\$F2,\$F4	\$F0=\$F2/\$F4	Single and double precision division
mov.s mov.d	mov.x \$F0,\$F2	\$F0<=\$F2	Single and double precision move
abs.s abs.d	abs.x \$F0,\$F2	\$F0=ABS(\$F2)	Single and double precision absolute value
neg.s neg.d	neg.x \$F0,\$F2	\$F0= - (\$F2)	Single and double precision absolute value
c.lt.s c.lt.d (eq.ne.le.gt.ge)	c.lt.x \$F0,\$F2	Temp=(\$F0<\$F2)	Single and double: compare \$F0 and \$F2 <,<=,>,>=
mtc1 (mfc1)	mtc1 \$1,\$F2	\$F2=\$1	Data from gen.reg. to C1 reg. (no conversion) (and viceversa)
branch on false	bolf label	If (Temp == false) go to label	Temp is 'Condition-Code'
branch on true	bolt label	If (Temp == true) go to label	Temp is 'Condition-Code'
load floating point (32bit)	lwc1 \$F0,0(\$1)	\$F0<=Memory[\$1]	
store floating point (32bit)	swc1 \$F0,0(\$1)	Memory[\$1]<=\$F0	
convert single into double	cvt.d.s \$F0,\$F2	\$F0=(double)\$F2	Also cvt.s.d (viceversa)
convert single into integer	cvt.w.s \$F1,\$F0	\$F1=(int)\$F0	Also cvt.s.w (viceversa)

**Register Usage**

Name	Register Num.	Usage	Name	Register Num.	Usage	Name	Usage
\$zero	0	The constant value 0	\$v0-\$v1	2-3	Results	\$f0, \$f1, ..., \$f31	Single precision floating point registers
\$s0-\$s7	16-23	Saved	\$fp, \$sp	30,29	frame pointer, stack pointer	\$f0, \$f2, ..., \$f30	Double precision floating point registers
\$t0-\$t9	8-15,24-25	Temporaires	\$ra, \$gp	31,28	return address, global pointer		
\$a0-\$a3	4-7	Arguments	\$k0-\$k1	26,27	Kernel usage		

**System calls**

Service Name	Service Num. (\$v0)	INPUT Arguments	OUTPUT Arguments
print_int	1	\$a0=integer to print	---
print_float	2	\$f12=float to print	---
print_string	4	\$a0=address of ASCII string to print	---
Sbrk	9	\$a0=Number of bytes to be allocated	\$v0=pointer to the allocated memory

- 1) Normalizzando 0.045 si ottiene:  $1.44 \cdot 2^{-5}$ . L' "uno" iniziale non viene rappresentato nel formato IEEE-754. Successivamente si puo' ricavare la rappresentazione binaria di 0.44 con 23 bit moltiplicando per 2 e ricavando la n-esima cifra piu' significativa della mantissa M:

$M = 0111\ 0000\ 1010\ 0011\ 1101\ 011$

(es.  $0.44 \cdot 2 = 0.88 \rightarrow 0$ ,  $0.88 \cdot 2 = 1.76 \rightarrow 1$ ,  $0.76 \cdot 2 = 1.52 \rightarrow 1$ ,  $0.52 \cdot 2 = 1.04 \rightarrow 1$ ,  $0.04 \cdot 2 = 0.08 \rightarrow 0$ , ...)

Per l'esponente, ricordando che nel caso di singola precisione il valore della polarizzazione e' 127:

$E = -5 + 127 = 122$  ovvero 01111010

Osservando che le cifre binarie di M si ripetono dalla 21-esima cifra e in particolare la 24-esima cifra e' un 1 si desume che la rappresentazione piu' vicina e' quella superiore. Quindi la rappresentazione cercata e':  
0 0111 1010 0111 0000 1010 0011 1101 100

- 2) Una possibile soluzione e' riportata nel file ncond.s.

```
.data
f:      .float 0.0
A:      .float 1.0, 2.0, 3.0, 4.0, 5.0, 6.0, 7.0, 8.0, 9.0
ret:    .ascii "\n"
esi:    .ascii "esito="
nco:    .ascii " n.cond="

.text
.globl main
.extern exp

# $a0=**T
# $a1=n
# $a2=*nc
# $v0=r
# $t0=j
# $t1=k
num_cond:
    addi    $sp, $sp, -24 # spazio per $a0,$v0,$t0,$t1
    sw      $ra, 0($sp) # salva $ra perche' usa altra f.
    sw      $fp, 4($sp) # salva $fp perche' usa frame
    add     $fp, $sp, $0

    add     $v0, $0, $0
    addi    $t0, $0, 1 # j = 1

    addi    $t2, $0, 1 # costante f.p. 1
    mtcl   $t2, $f1
    cvt.s.w $f1, $f1
    addi    $t2, $0, 0 # costante f.p. 0
    mtcl   $t2, $f0 # la codifica di 0.0 e' 0...0

while_ini:
    slt    $t2, $a1, $t0 # j>?n
    bne    $t2, $0, while_end # se si while_end

    add     $t1, $0, $0 # k = 0
for_ini:
    slt    $t2, $t1, $a1 # k<?n
    beq    $t2, $0, for_end # se no for_end

    addi    $t3, $t0, -1 # j -1
    add     $t2, $t3, $t3
    add     $t2, $t2, $t3 # (j-1)*3
    add     $t2, $t2, $t1 # (j-1)*3+k
    sll    $t2, $t2, 2 # *4
    add     $t2, $t2, $a0 # &T[j-1][k]
    lwc1   $f2, 0($t2) # T[j-1][k]

    c.eq.s $f2, $f0 # == 0.0
    bclt   ramo_else

    sw      $a0, 8($fp) # salva $a0
    sw      $v0, 12($fp) # salva $v0
    sw      $t0, 16($fp) # salva $t0
    sw      $t1, 20($fp) # salva $t1

    neg.s   $f2, $f2 # cambio segno
    mfcl   $a0, $f2
    jal    exp # $v0=exp($a0)
    mtcl   $v0, $f2

    lw      $t1, 20($fp) # salva $t1
    lw      $t0, 16($fp) # salva $t0
    lw      $v0, 12($fp) # ripristina
    lw      $a0, 8($fp) # ripristina

    div.s   $f2, $f1, $f2 # 1/sqr(-
T[[]])

    lwc1    $f3, 0($a2) # *nc
    add.s   $f3, $f3, $f2 # +=
    swc1    $f3, 0($a2) # *nc=
    j       if_end

ramo_else:
    addi    $v0, $0, 1 # r = 1

if_end:
    addi    $t1, $t1, 1 # ++k
    j       for_ini

for_end:
    addi    $t0, $t0, 1 # ++j
    j       while_ini

while_end:
    lw      $ra, 0($fp)
    lw      $fp, 4($fp)
    addi    $sp, $sp, 24
    jr      $ra # ritorna $v0

main:
    la     $a0, A # param.1
    addi   $a1, $0, 3 # param.2
    la     $a2, f # param.3

    jal    num_cond

    add     $s0, $v0, $0 # salva ec

    la     $a0, esi #print "esi..."
    addi   $v0, $0, 4
    syscall

    add     $a0, $s0, $0 # ripristina
ec:
    addi   $v0, $0, 1
    syscall
    la     $a0, nco # print

"nco..."
    addi   $v0, $0, 4
    syscall
    la     $a0, f # print f
    lwc1   $f12, 0($a0)
    addi   $v0, $0, 2
    syscall
    la     $a0, ret # print ret
    addi   $v0, $0, 4
    syscall
    addi   $v0, $0, 10 # exit
    syscall
```