

1) [30/40] Trovare il codice assembly MIPS corrispondente del seguente programma (utilizzando solo e unicamente istruzioni dalla tabella sottostante e rispettando le convenzioni di utilizzazione dei registri dell'assembly MIPS riportate qua sotto, per riferimento).

```

/* Finds the index of element having max. absolute value.
** Jack Dongarra, linpack, 3/11/78.
*/
static int idamax(int n, float *dx,int incx)
{
    float dmax;
    int i, ix, itemp;

    if (n < 1)
        return(-1);
    if (n ==1 )
        return(0);
    if(incx != 1)
    {
        /* code for increment not equal to 1 */

        ix = 1;
        dmax = fabs((double)dx[0]);
        ix = ix + incx;
        for (i = 1; i < n; i++)
        {
            if(fabs((double)dx[ix]) > dmax)
            {
                itemp = i;
                dmax = fabs((double)dx[ix]);
            }
            ix = ix + incx;
        }
    }
    else
    {
        /* code for increment equal to 1 */

        itemp = 0;
        dmax = fabs((double)dx[0]);
        for (i = 1; i < n; i++)
            if(fabs((double)dx[i]) > dmax)
            {
                itemp = i;
                dmax = fabs((double)dx[i]);
            }
    }
    return (itemp);
}

int main()
{
    float a[100];
    int l, k, n, lda;

    l = idamax(n-k,&a[lda*k+k],1) + k;
}

```

MIPS instructions

Instruction	Example	Meaning	Comments
add	add \$1,\$2,\$3	\$1 = \$2 + \$3	3 operands; exception possible
subtract	sub \$1,\$2,\$3	\$1 = \$2 - \$3	3 operands; exception possible
add immediate	addi \$1,\$2,100	\$1 = \$2 + 100	+ constant; exception possible
subtract immediate	subi \$1,\$2,100	\$1 = \$2 - 100	- constant; exception possible
multiplication	mult \$1, \$2	Hi,Lo= \$1 x \$2	64-bit Signed Product ; result in Hi,Lo
division	div \$1, \$2	Hi= \$1 % \$2, Lo = \$1 / \$2	Signed division
move from Hi	mfhi \$1	\$1 = Hi	Create copy of Hi
move from Lo	mflo \$1	\$1 = Lo	Create copy of Lo
and	and \$1,\$2,\$3	\$1 = \$2 & \$3	3 register operands; Logical AND
or	or \$1,\$2,\$3	\$1 = \$2 \$3	3 register operands; Logical OR
nor	nor \$1,\$2,\$3	\$1 = ~(\$2 \$3)	3 register operands; Logical NOR
xor	xor \$1,\$2,\$3	\$1 = \$2 ^ \$3	3 register operands; Logical XOR
and immediate	andi \$1,\$2,100	\$1 = \$2 & 100	Logical AND register, constant
or immediate	ori \$1,\$2,100	\$1 = \$2 100	Logical OR register, constant
xor immediate	xori \$1,\$2,100	\$1 = \$2 ^ 100	Logical XOR register, constant
shift left logical	sll \$1,\$2,10	\$1 = \$2 << 10	Shift left by constant
shift right logical	srl \$1,\$2,10	\$1 = \$2 >> 10	Shift right by constant
load word	lw \$1,100(\$2)	\$1 = Memory[\$2+100]	Data from memory to register
load byte	lb \$1,100(\$2)	\$1 = Memory[\$2+100]	Data from memory to register
load byte unsigned	lbu \$1,100(\$2)	\$1 = Memory[\$2+100]	Data from mem. to reg.; no sign extension
store word	sw \$1,100(\$2)	Memory[\$2+100] = \$1	Data from register to memory
store byte	sb \$1,100(\$2)	Memory[\$2+100] = \$1	Data from register to memory
load address	la \$1,var	\$1 = &var	Load variable address
branch on equal	beq \$1,\$2,100	if (\$1 == \$2) go to PC+4+100	Equal test; PC relative branch
branch on not equal	bne \$1,\$2,100	if (\$1 != \$2) go to PC+4+100	Not equal test; PC relative
set on less than	slt \$1,\$2,\$3	if (\$2 < \$3) \$1 = 1; else \$1 = 0	Compare less than; 2's complement
set on less than immediate	slti \$1,\$2,100	if (\$2 < 100) \$1 = 1; else \$1 = 0	Compare < constant; 2's complement
set on less than unsigned	sltu \$1,\$2,\$3	if (\$2 < \$3) \$1 = 1; else \$1 = 0	Compare less than; natural number
set on less than imm. unsigned	sltiu \$1,\$2,100	if (\$2 < 100) \$1 = 1; else \$1 = 0	Compare constant; natural number
jump	j 10000	go to 10000	Jump to target address
jump register	jr \$31	go to \$31	For switch, procedure return
jump and link	jal 10000	\$31 = PC + 4; go to 10000	For procedure call
add.s add.d	add.x \$F0,\$2,\$F4	\$F0=\$F2+\$F4	Single and double precision add
sub.s sub.d	add.x \$F0,\$2,\$F4	\$F0=\$F2-\$F4	Single and double precision subtraction
mul.s mul.d	mul.x \$F0,\$2,\$F4	\$F0=\$F2*\$F4	Single and double precision multiplication
div.s div.d	div.x \$F0,\$2,\$F4	\$F0=\$F2/\$F4	Single and double precision division
mov.s mov.d	mov.x \$F0,\$F2	\$F0←\$F2	Single and double precision move
abs.s abs.d	abs.x \$F0,\$F2	\$F0=ABS(\$F2)	Single and double precision absolute value
neg.s neg.d	neg.x \$F0,\$F2	\$F0 = -(\$F2)	Single and double precision absolute value
c.lts c.lt.d (eq.ne,le.gt,ge)	c.lt.l.x \$F0,\$F2	Temp=(S0<Sf2)	Single and double: compare \$f0 and \$f2 <=,!=,<<,>>
mtcl (mfcl)	mtcl \$1,\$F2	\$F2-\$f1	Data from gen.reg. to C1 reg. (no conversion) (and viceversa)
branch on false	bclf label	If (Temp == false) go to label	Temp is 'Condition-Code'
branch on true	bclt label	If (Temp == true) go to label	Temp is 'Condition-Code'
load floating point (32bit)	lwc1 \$F0,0(\$1)	\$F0←Memory[\$1]	
store floating point (32bit)	swc1 \$F0,0(\$1)	Memory[\$1]←\$F0	
convert single into double	cvt.d.s \$F0,\$F2	\$F0=(double)\$F2	Also cvt.s.d (viceversa)
convert single into integer	cvt.w.s \$f1,\$F0	\$f1=(int)\$f0	Also cvt.s.w (viceversa)

Register Usage

Name	Register Num.	Usage
\$zero	0	The constant value 0
\$s0-\$s7	16-23	Saved
\$t0-\$t9	8-15,24-25	Temporaries
\$a0-\$a3	4-7	Arguments

Name	Register Num.	Usage
\$v0-\$v1	2-3	Results
\$fp, \$sp	30,29	frame pointer, stack pointer
\$ra, \$gp	31,28	return address, global pointer
\$k0-\$k1	26,27	Kernel usage

Name	Usage
\$f0, \$f1, ..., \$f31	Single precision floating point registers
\$f0, \$f2, ..., \$f30	Double precision floating point registers

System calls

Service Name	Service Num. (\$v0)	INPUT Arguments	OUTPUT Arguments
print_int	1	\$a0=integer to print	---
print_float	2	\$f12=float to print	---
print_string	4	\$a0=address of ASCII string to print	---
Sbrk	9	\$a0=Number of bytes to be allocated	\$v0=pointer to the allocated memory

2) [10/40] Si consideri una cache di dimensione 64B e a 4 vie di tipo write-back. La dimensione del blocco e' 8 byte, il tempo di accesso alla cache e' 4 ns e la penalita' in caso di miss e' pari a 40 ns, la politica di rimpiazzamento e' LRU. Il processore effettua i seguenti accessi in cache, ad indirizzi al byte: 7773, 7756, 7777, 7790, 7774, 7789, 7773, 7788, 7779, 5390, 7777, 7772, 7703, 5270, 7760, 7775, 7787, 7771, 7707, 5378, 7712.

Tali accessi sono alternativamente letture e scritture. Per la sequenza data, ricavare il tempo medio di accesso alla cache, riportare i tag contenuti in cache al termine e la lista dei blocchi (ovvero il loro indirizzo) via via eliminati durante il rimpiazzamento ed inoltre in corrispondenza di quale riferimento il blocco e' eliminato.